

Targeted Search: Reducing the Time and Cost for Searching for Objects in Multiple Server Networks

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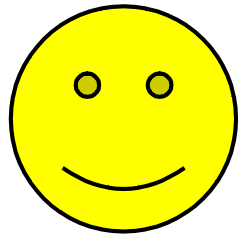
Topics

❖ This talk is about


1) Traffic

2) Flooding

Traffic and flooding!



Topics continued

- ❖ Introduction
- ❖ Targeted search method  **New**
- ❖ Performance evaluation
- ❖ Ditella – Targeted search in Gnutella
- ❖ Summary and future work

Introduction

- ❖ **Question: How to efficiently find an object in multiple servers?**
 - For example, in a P2P network
- ❖ **Efficient = fast with low overhead**
 - Overhead is query traffic on a P2P network
- ❖ **We have a trade-off between...**
 - Time to complete a search
 - Amount of overhead traffic



Much of the traffic on the Internet is P2P search queries

Introduction continued

❖ P2P networks are an example of multiple server networks

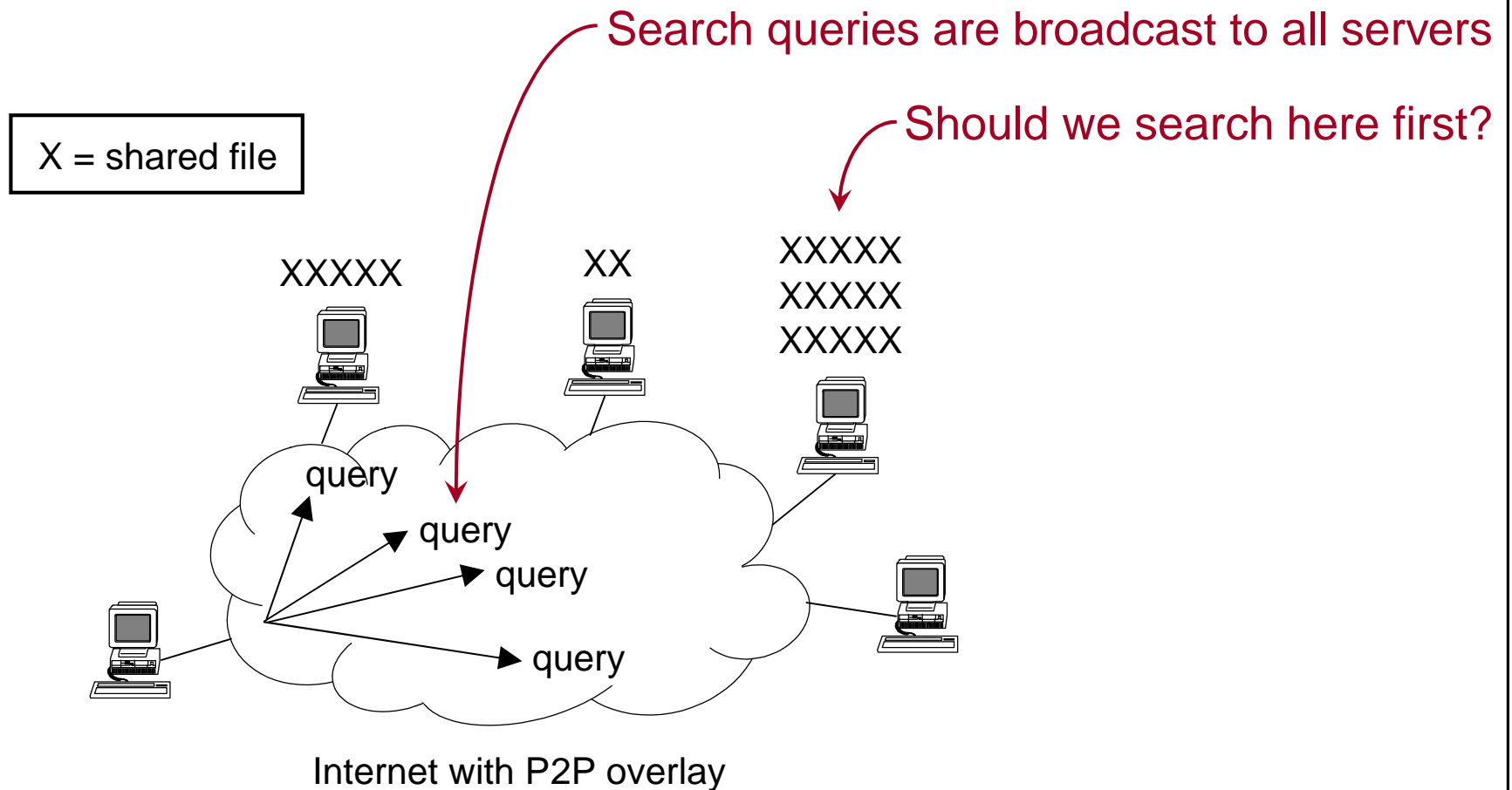
- Used for file sharing

❖ P2P search methods

- Napster
 - Single master directory
- Gnutella
 - No master directory – broadcast a query to all “servents” (nodes)
- Kazaa
 - Hierarchical directory structure (many local directories)
- Chord
 - Hash of file name determines where a file is stored

Introduction continued

- ❖ **Power-law distribution of files is a reasonable assumption**
 - Few servers share most of files



Introduction continued

❖ Our hypothesis...

Individual nodes can “learn” the distribution of objects in servers and use this learned knowledge to improve searching.

❖ Search first in servers with the most files

- Highest probability of finding file

❖ May not have to broadcast to all servers

- Thus, reduce the amount of query traffic

Introduction continued

- ❖ Can prove that an optimal search exists
- ❖ Determine optimal search method for given time constraint k
- ❖ Have N servers and uniform probability of file being in a server
- ❖ Check at each step $s(1), s(2), \dots, s(k)$ such that $s(1) + s(2) + \dots + s(k) = N$
- ❖ The mean cost is

$$E[C] = N/2 + (s(1)^2 + s(2)^2 + \dots + s(k)^2)/2N$$

- ❖ The strategy that minimizes $E[C]$ is $s(1) = s(2) = \dots = s(k) = N/k$ where

$$E[C] = N(k+1)/2k$$



For non-uniform distribution we show that weighted search is optimal (see paper for proofs)

Targeted search method

New

- ❖ A method to target high-probability nodes for a direct query
- ❖ Each node maintains a frequency list
 - Tuples of `<server_id, hit_count>`
- ❖ Frequency list is sorted by previous successes
 - Sort variable is `hit_count`



Do a direct search to the N_{top} highest ranked nodes and if the file is not found, then broadcast a query to all nodes

Targeted search method continued

❖ The algorithm

- Node searching for a file in N servers

Step 1: Send a direct query iteratively to the servers in the frequency list starting with the first listed server. This step terminates when N_{top} servers have been queried or the object has been found.

Step 2: If Step 1 did not find the object then broadcast a query to all N servers.

Step 3: If Step 1 or Step 2 found the object in a server then download the object and update the frequency list.

Performance evaluation

❖ Performance evaluation using analysis and simulation

- Analysis with theoretical distribution of files
- Simulation with traced (empirical) distribution of files

❖ Control variable:

- Distribution of files in nodes

❖ Response variables:

- Time = proportional to number of query steps
- Cost = number of nodes queried



We seek closed form expressions for mean time and cost

Performance evaluation continued

- ❖ Define a peaked distribution for $\Pr[\text{file in server } i]$
 - N is number of servers
 - K is a “tuning knob”

$$\Pr[\text{file in server } i] = \begin{cases} \frac{K}{N+K-1} & i=1 \\ \frac{1}{N+K-1} & i=2,3,\dots,N \end{cases}$$

- ❖ $\Pr[\text{file in server } i = 1]$ is K times greater than $\Pr[\text{file in server } i > 1]$
- ❖ As K increases $\Pr[\text{file in server } i = 1]$ goes to 1
- ❖ As K increases and $\Pr[\text{file in server } i > 1]$ goes to 0

Performance evaluation continued

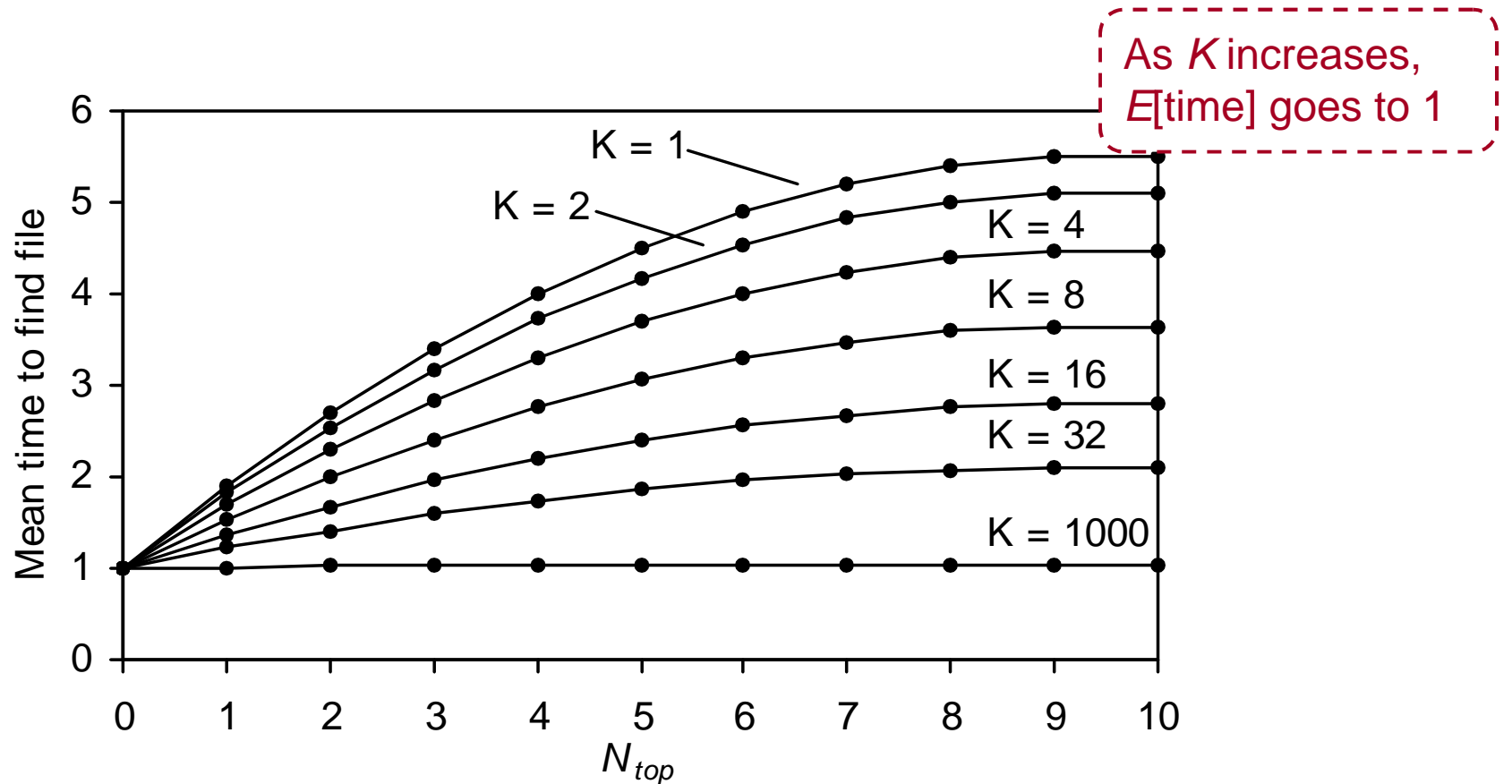
- ❖ We wish to find an expression for
 - The mean time to find an object ($E[\text{cost}]$)
 - The mean cost to find an object ($E[\text{time}]$)
- ❖ For files distributed with the peaked distribution we find

$$E[\text{time}] = \frac{2N_{top}N - N_{top}^2 - N_{top} + 2 \cdot N + 2K - 2}{2(N + K - 1)}$$

$$E[\text{cost}] = \frac{2N^2 - N_{top}^2 + N_{top} + 2K - 2}{2(N + K - 1)}$$

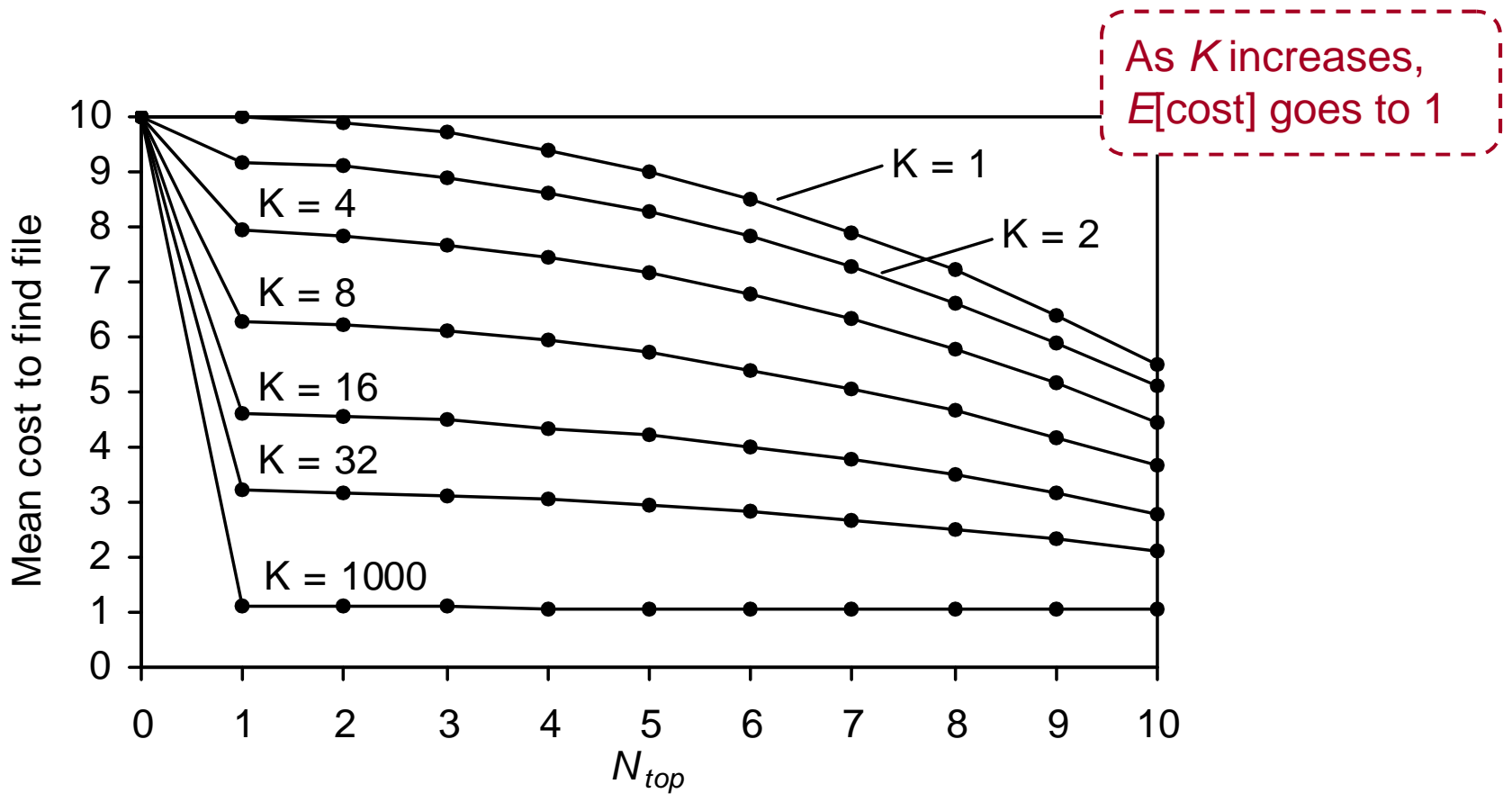
Performance evaluation continued

❖ Analytical mean time results for $N = 10$



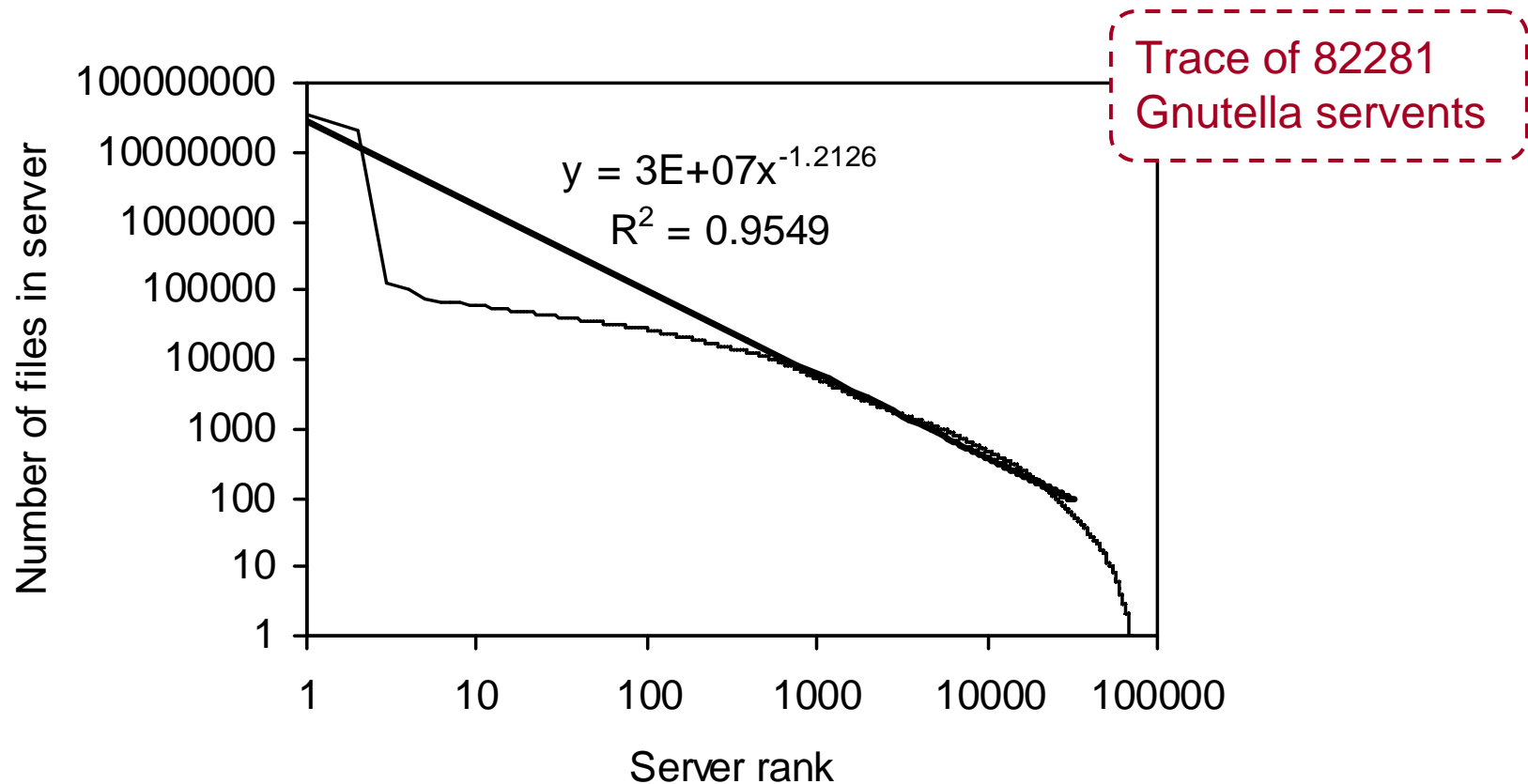
Performance evaluation continued

❖ Analytical mean cost results for $N = 10$



Performance evaluation continued

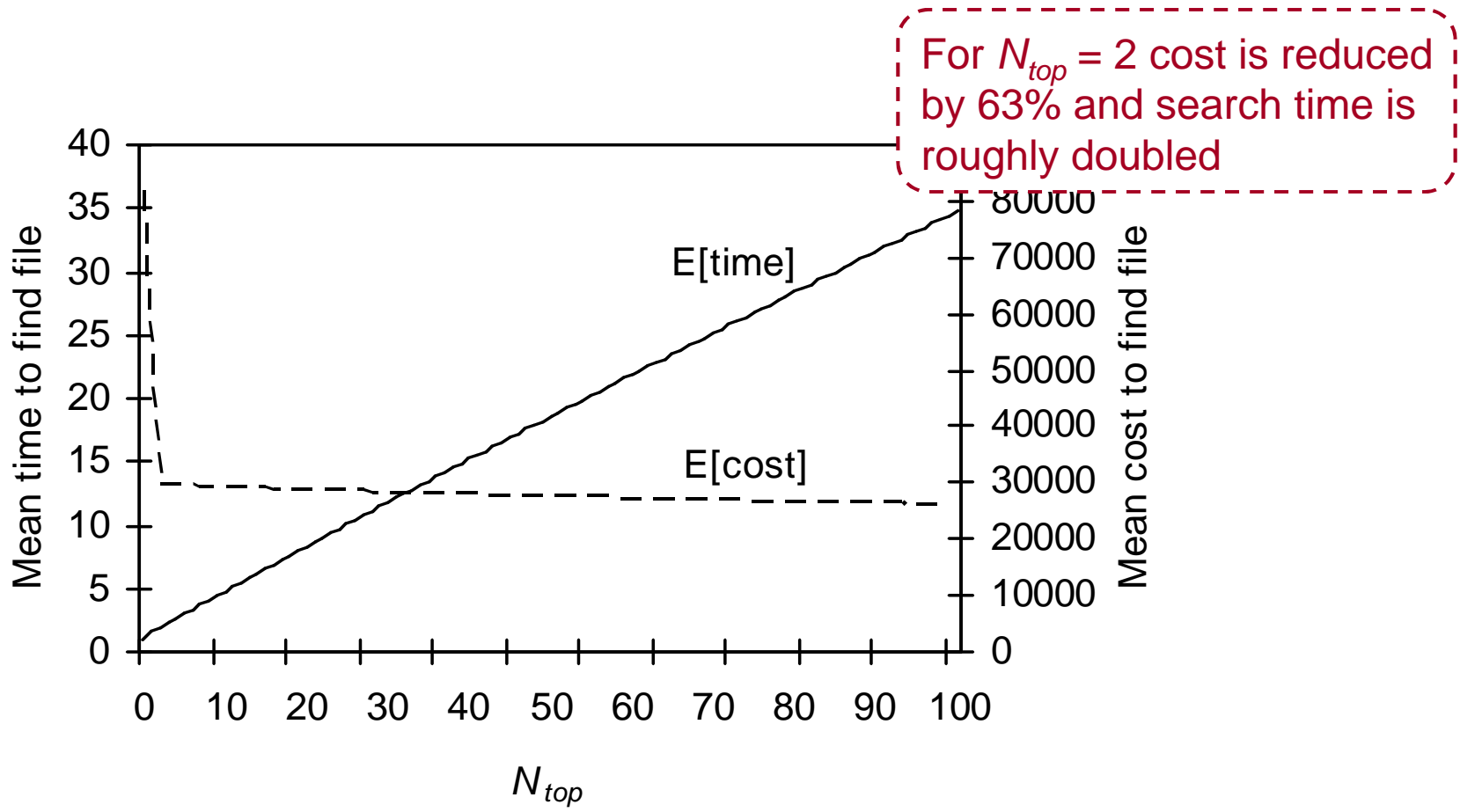
- ❖ **Simulation evaluation using Gribble et al. file trace for input**
 - See Power Law fit on log-log graph



S. Saroiu, P. Gummadi, and S. Gribble, "A Measurement Study of Peer-to-Peer File Sharing Systems," *Proceedings of SPIE*, pp. 156-170, 2001.

Performance evaluation continued

❖ Simulation mean time and cost results



Performance evaluation continued

❖ See paper for...

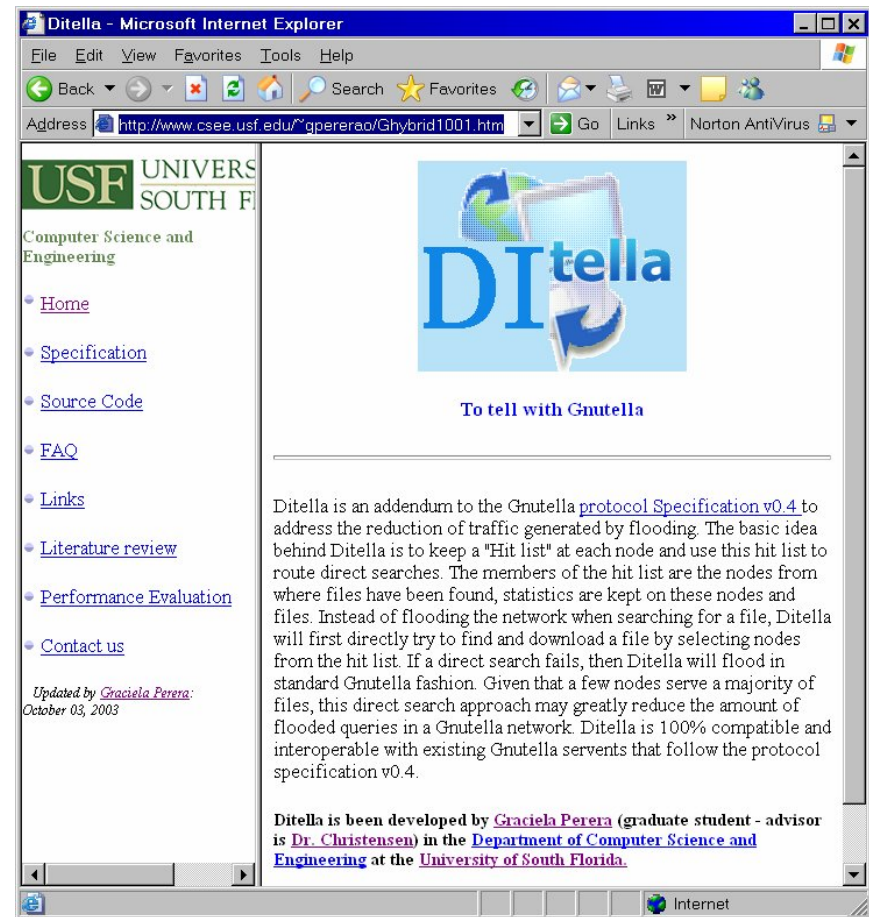
- Comparison to random walk P2P search
- Evaluation of convergence of frequency list to distribution

Ditella – targeted search in Gnutella

❖ Implemented targeted search

- In a Gnutella servant
- For Microsoft Windows

❖ Available at project website



<http://www.csee.usf.edu/~gpererao/Ghybrid1001.htm>

Summary and future work

- ❖ Distribution of files to servers is often power law
- ❖ Can exploit this to reduce search time and cost
- ❖ Our new *targeted search* is: **query first the most popular nodes**
 - Performance evaluation with analysis and simulation
- ❖ Implemented targeted search in Gnutella
 - We call this Ditella
- ❖ Future work a better (more usable and portable) implementation
 - Goal is widespread adoption and deployment

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